

# Minimum Circuit Size Problem

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# Introduction

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- 2 The Minimum Circuit Size Problem **MCSP** is an example of a problem of unknown complexity.

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## NP-completeness

A problem is **NP-Complete** if every problem in NP reduces to it.  
If any **NP-Complete** problem is in P then  $P = NP$

## Circuit

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## MCSP

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MCSP is suspected to be an NP-Intermediate problem. This means it is not solvable in polynomial time and it is in NP, but it is also not NP-Hard.

## RP, co-RP, ZPP

- 1 **RP** is the class of algorithms that run in polynomial time with small probability of rejecting instead of accepting.
- 2 **co-RP** is the class of algorithms that run in polynomial time with small probability of accepting instead of rejecting.
- 3 **ZPP** is the class of algorithms that run with no error in expected polynomial time.

## Oracle classes

Class  $C^A$  contains languages which can be solved with an algorithm using computational power  $C$  and a subroutine to decide the problem  $A$ .

# Graph Automorphism

**Graph Automorphism** Problem: Given Graph  $G$  is there a permutation  $\sigma$  such that edge  $(u, v) \in G$  if and only if  $(\sigma(u), \sigma(v)) \in \sigma(G)$ .

This is another problem in NP which is not known to be in P and not known to be NP-Complete

## Open problem

Is  $GA \in ZPP^{MCSP}$ ?

Approach: prove  $GA \in RP^{MCSP}$  and  $GA \in co - RP^{MCSP}$ .

To prove  $GA \in RP^{MCSP}$ , attempt to follow the technique used by Allender to prove  $GI$  (Graph Isomorphism)  $\in RP^{MCSP}$  [?]

Given  $GA \in RP^{MCSP}$ , if  $GA \in ZPP^{MCSP}$  then  $GnA \in ZPP^{MCSP}$ . Thus it is likely difficult to prove  $GA \in co - RP^{MCSP}$

## Interactive Proofs

The **prover** wants to persuade the **verifier** that something holds (US beer vs water).

## Statistical Zero Knowledge

Interactive proof where the verifier learns statistically nothing about the solution (will not learn how to tell the drinks apart).

# Kolmogorov Complexity

**Kolmogorov complexity** of a string  $x$  is the length of the shortest program that outputs  $x$ .

A string is **Kolmogorov random** if its complexity is at least its length.

## Open Question (Allender)

Is there a significant subclass of SZK that reduces to Kolmogorov random strings?

# Conclusion

- 1 MCSP is a difficult problem. It is been studied for decades and still defies provable classification with relation to  $P$  and being  $NP - Hard$ .
- 2 Through determining it's relative complexity compared to other problems of unknown complexity, the difficulty of MCSP can be more accurately determined.
- 3 The goal of this summer is to examine several open problems focusing on various reductions between MCSP and other possibly NP-Intermediate problems to provide evidence for future classification of MCSP.

# TODO bibliography

Thank you for your attention.